

# On the analysis of Internet Paths with DisNETPerf, a Distributed Internet Paths Performance Analyzer

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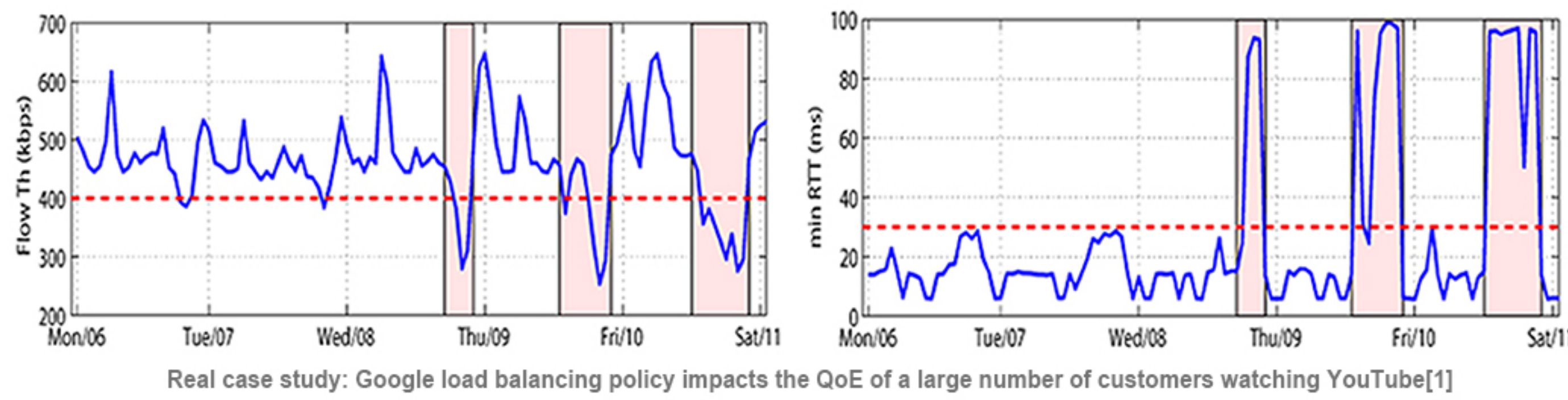
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## Why DisNETPerf?

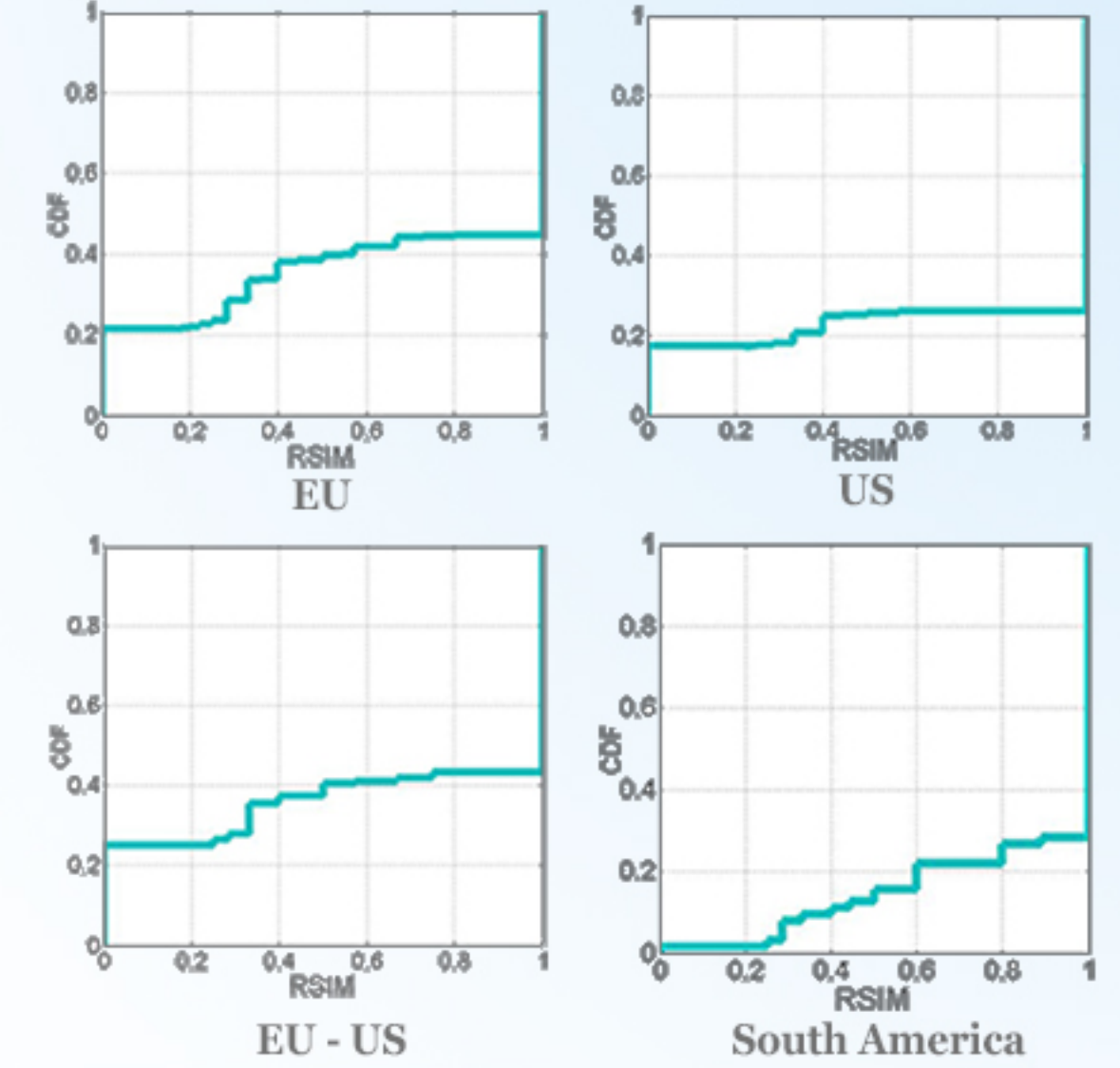
- Load-balancing: user-requests redirected to servers at hundreds of milliseconds
- Performance-monitoring of reverse paths (uncontrolled server→customer) challenging
- Reverse traceroute[2] relies on IP spoofing and IP Record Option: security concerns



Real case study: Google load balancing policy impacts the QoE of a large number of customers watching YouTube[1]

## Path asymmetry study

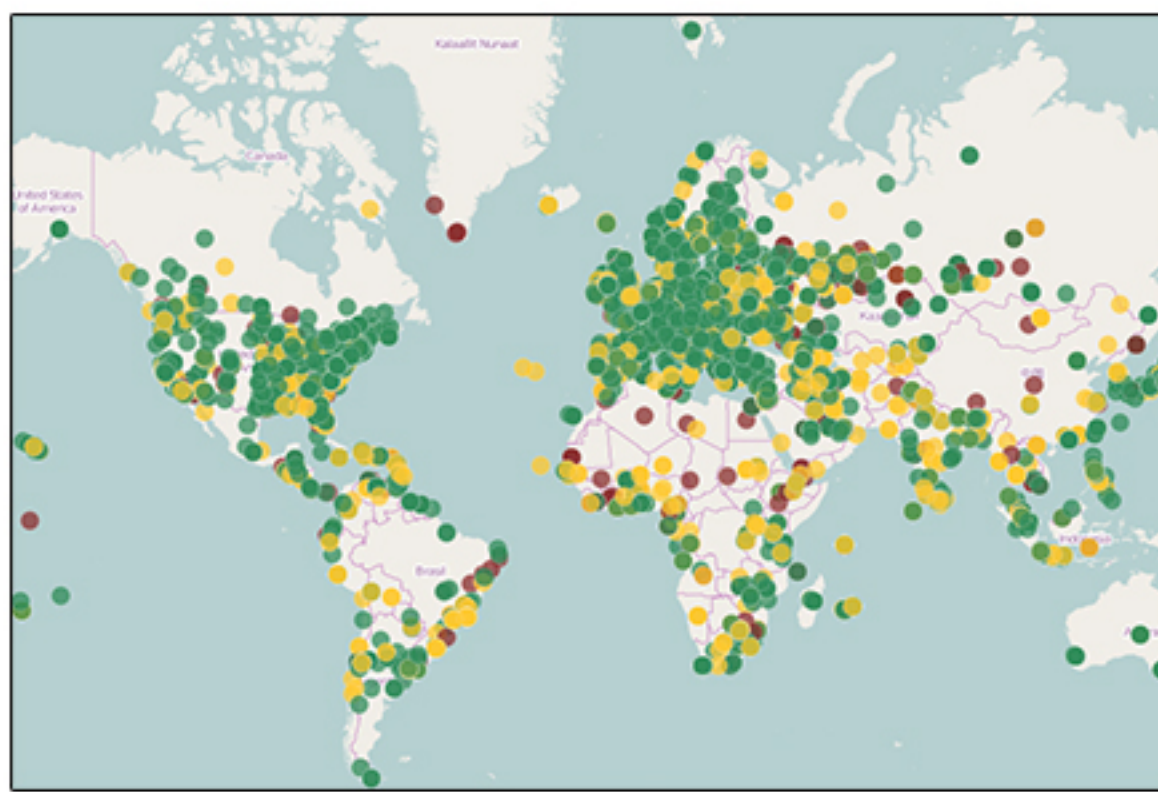
- Path asymmetry study between RIPE Atlas boxes in different regions of the world
- Asymmetric path:  $A \rightarrow B$  different from  $B \rightarrow A$
- Metric: *Route Similarity* index (RSIM) - ratio between #shared links and #total links
- Path asymmetry indicated by  $RSIM(A \rightarrow B, \text{inverse}(B \rightarrow A))$
- More asymmetric paths in larger zones
- Path asymmetry non negligible phenomenon!



## Overview

- DisNETPerf relies on RIPE Atlas (RA) distributed measurements
- Given: content server  $IP_s$ , destination  $IP_d$ , locate  $IP_c$  - closest RA probe to  $IP_s$
- Combined topology- and delay-based distance: first locate candidate RA probes by AS and then elect  $IP_c$  by propagation delay
- Periodically run *traceroute* measurements from  $IP_c$  to  $IP_d$
- Data can be used to troubleshoot paths from the content server (mimicked by  $IP_c$ ) to the target customer

Distribution of RA probes



## Probe selection: 2 approaches [3]

### Smallest latency (SL) approach

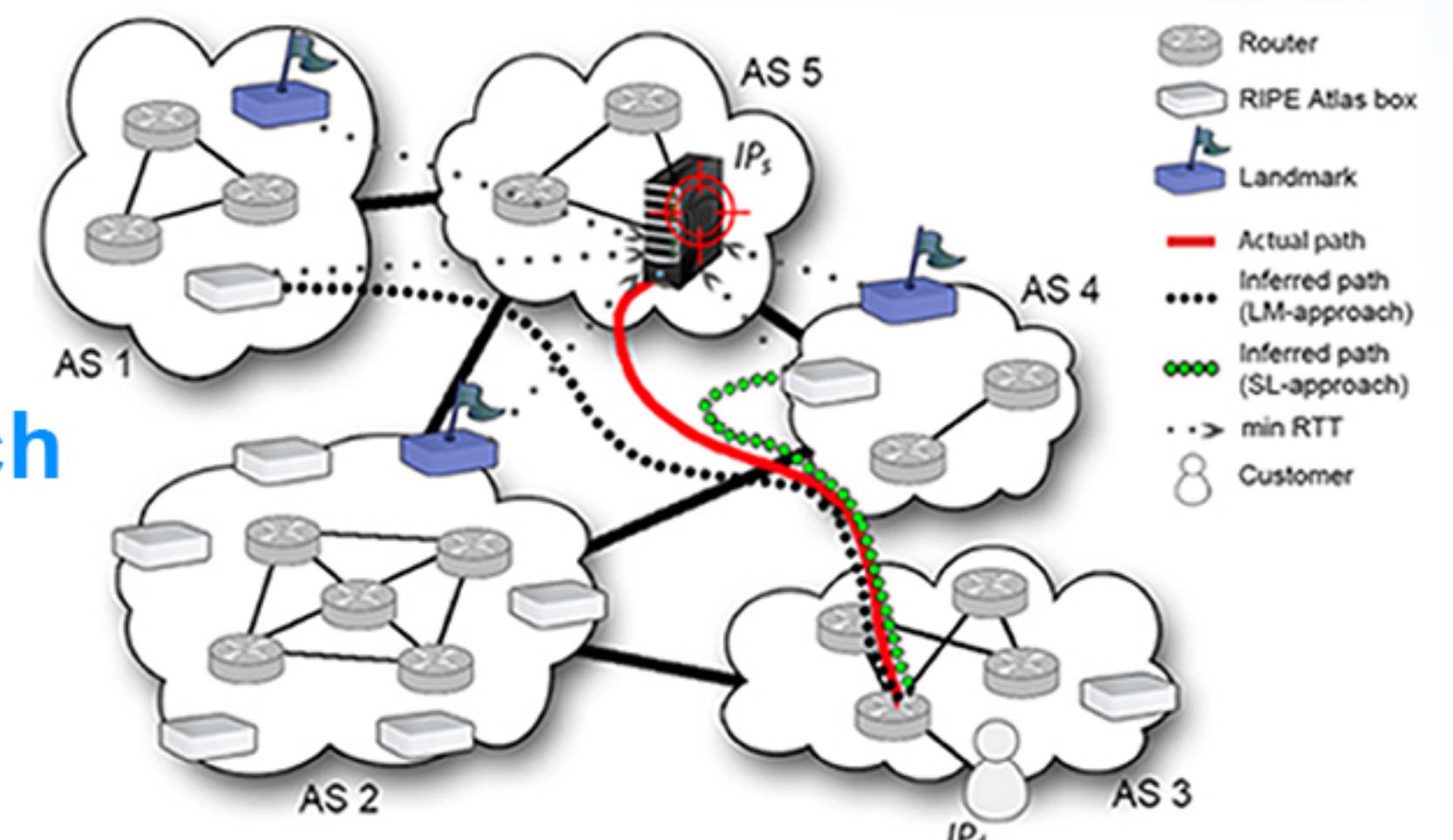
- Launch *ping* measurements from candidate RA probes to  $IP_s$
- Probe with smallest minimum RTT elected as  $IP_c$

### Landmark (LM) approach

- Use some candidate RA probes as landmarks for distance reference
- Launch *ping* measurements from landmarks to  $IP_s$  and all other candidate RA probes
- Elect as  $IP_c$  probe with smallest normalized distance to  $IP_s$ :

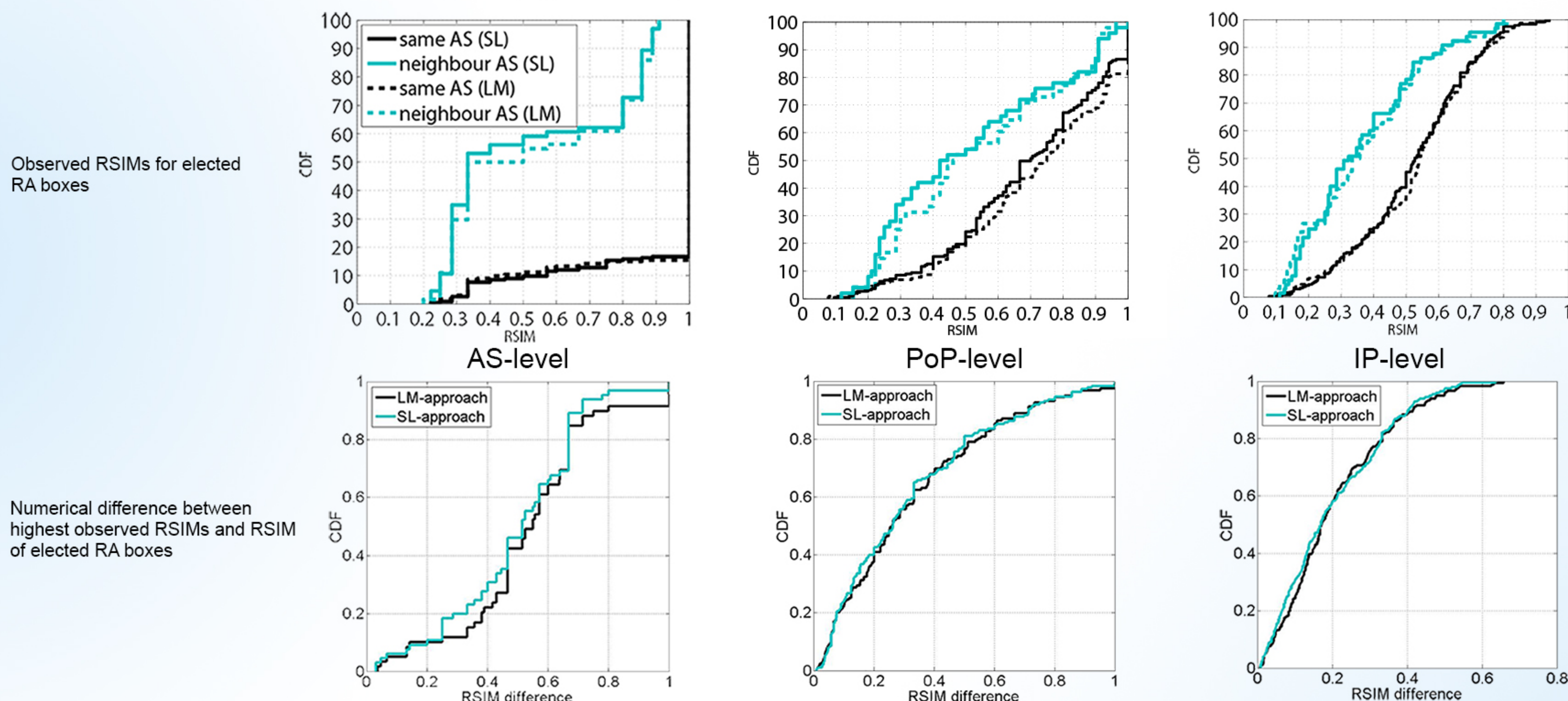
$$D_{ij} = \frac{1}{K} \sum_{l=1}^K |d_{il} - d_{jl}|$$

where  $K$  = #landmarks providing a RTT for both  $IP_i$  and  $IP_j$ ,  
 $d_{il}$  = min RTT between  $IP_i$  and landmark  $l$



## Evaluation

- $IP_c$  is a good probe w.r.t.  $IP_s$  and  $IP_d$  if the path from  $IP_c$  to  $IP_d$  is highly similar to the path from  $IP_s$  to  $IP_d$
- Route Similarity* index -  $RSIM(IP_c, IP_s, IP_d) = \frac{2 \times C_{links}(IP_c, IP_s, IP_d)}{T_{links}(IP_c, IP_s, IP_d)}$  where  $C_{links}$  = #common links,  $T_{links}$  = #total links
- Ground-truth: RA probes as  $IP_s$  and  $IP_d$  to compute real server→customer paths (300 tests)
- RSIM computed for different link definitions (AS, PoP, and IP level)
- AS-level measurements more relevant (troubleshooting faulty ASes) → high RSIM when RA probe located in same AS
- 2 approaches very accurate: select the best RA probe (highest RSIM) for more than 80% of tests



- Ongoing:** combined analysis of passive (@access ISP) and active measurements (DisNETPerf) for Youtube analysis

## References

[1] P. Casas, A. D'Alconzo, P. Fiadino, A. Bär, A. Finamore, and T. Zseby, "When YouTube does not work. analysis of QoE-relevant degradation in Google CDN traffic," *IEEE Transactions on Network and Service Management*, vol. 11, no. 4, pp. 441–457, December 2014.

[2] E. Katz-Bassett, H. Madhyastha, V. Adhikari, C. Scott, J. Sherry, P. van Wesep, A. Krishnamurthy, and T. Anderson, "Reverse traceroute," in *Proc. USENIX Symposium on Networked Systems Design and Implementations (NSDI)*, June 2010.

[3] S. Wassermann, P. Casas, B. Donnet, "Towards DisNETPerf: a Distributed Internet Paths Performance Analyzer", in *Proc. ACM CoNEXT Student Workshop*, December 2015